



**AN EMPIRICAL ANALYSIS AND SIMULATION OF A COLLISION
AVOIDANCE ALGORITHM FOR SECURE DATA EXCHANGE
AND NON REDUNDANT PACKET TRANSMISSION**

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Abstract

The secure transmission of data in transit relies on both cryptography and authentication – on both the hiding or concealment of the data itself, and on ensuring that the computers at each end are the computers they say they are. In this paper a priority based collision avoidance algorithm for secure data exchange is proposed where the encryption and decryption process uses hash function and authentication is supported via username password. The algorithm provides the unique feature of collision and redundancy avoidance and also supports the priority based response from receiver of data or message. This paper simulates the error free, non redundant and collision free packet transmission in the network channel

Keywords: *Priority, collision, data security, authentication, cryptography.*



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1 INTRODUCTION

The important security risk is that information can be captured and read during its transmission. How do we protect this information from being read by intruders? The secure transmission of data in transit relies on both cryptography and authentication – on both the hiding or concealment of the data itself, and on ensuring that the computers at each end are the computers they say they are.

Cryptography, to most people, is concerned with keeping communications private. Encryption is the transformation of data into some unreadable form. Its purpose is to ensure privacy by keeping the information hidden from anyone for whom it is not intended. Decryption is the reverse of encryption; it is the transformation of encrypted data back into some intelligible form. Encryption and decryption require the use of some secret information, usually referred to as a key. The data to be encrypted is called as plain text. The encrypted data obtained as a result of encryption process is called as cipher text. Depending on the encryption mechanism used, the same key might be used for both encryption and decryption, while for other mechanisms, the keys used for encryption and decryption might be different. There are several ways of classifying cryptographic algorithms. In general they are categorized based on the number of keys that are employed for encryption and decryption, and further defined by their application and use as in [1]. The three types of algorithms are depicted as follows

1) *Secret Key Cryptography (SKC)*: Uses a single key for both encryption and decryption. The most common algorithms in use include Data Encryption Standard (DES), Advanced Encryption Standard (AES).

2) *Public Key Cryptography (PKC)*: Uses one key for encryption and another for decryption. RSA (Rivest, Shamir, Adleman) algorithm is an example.



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3) *Hash Functions*: Uses a mathematical transformation to Irreversibly "encrypt" information.
MD (Message Digest)

In this paper a priority based collision avoidance algorithm for secure data exchange is proposed where the encryption and decryption process uses hash function and authentication is supported via username password. The algorithm provides the unique feature of collision and redundancy avoidance and also supports the priority based response from receiver of data or message. Section II presents the detailed study of hash function. Section III describes the proposed algorithm for secure data exchange. Finally section IV concludes the paper and presents the future work directions.

2. Proposed Algorithm and Data Flow Diagram

In this section a priority based collision avoidance algorithm for secure data exchange is proposed which provides the security to the long message in terms of authentication and confidentiality so that the privacy of data does not breach and not any unauthorized user access the data. This algorithm avoids the collision and storage of repeated data when similar messages repeated over network at times and also provides the facility to receiver for responding according to the priority wise and separately to each sender very efficiently. The functional flow diagram and algorithm steps are described in figure 6 and 7 respectively

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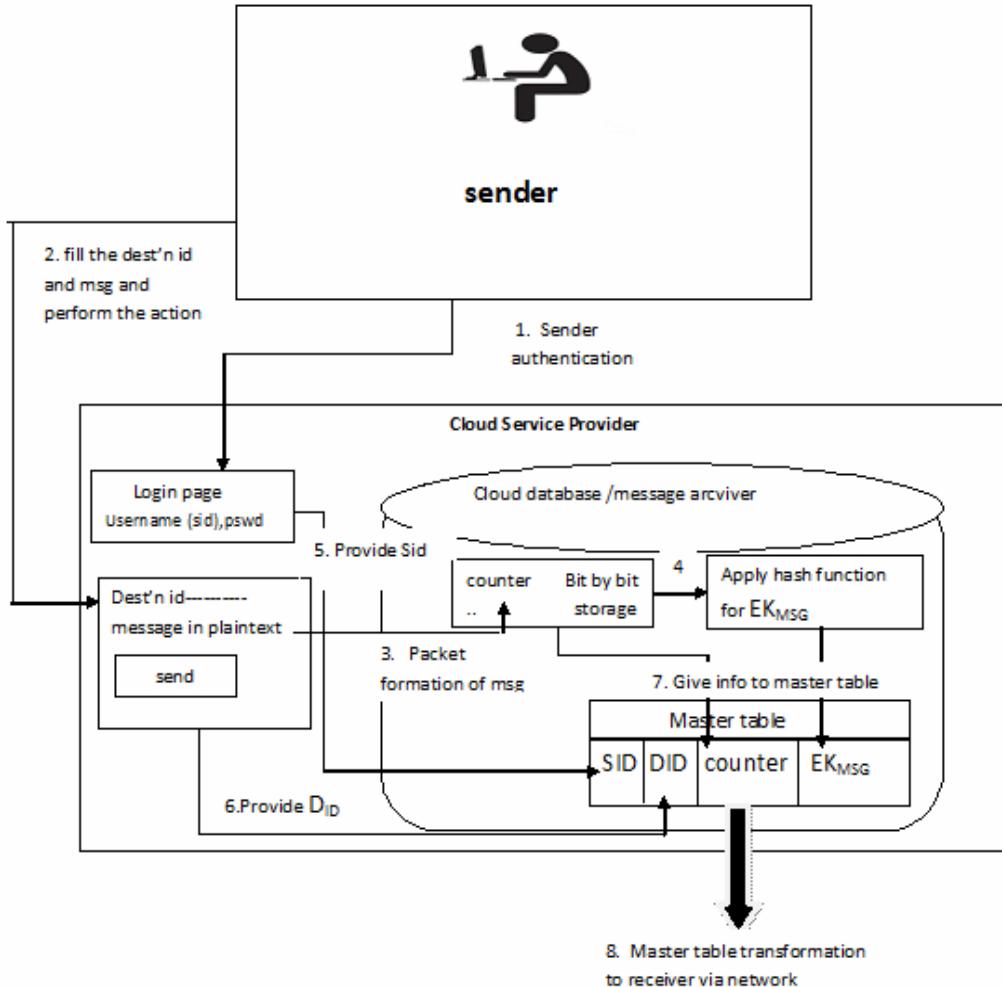


Figure 6. Functional flow diagram for sending encrypted message

Algorithm for sending encrypted message or data

- ### 1. Authentication of sender



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(a) Sender login to cloud server where:

$S_{ID} \leftarrow \text{Username}$ /* set username of sender as sender id S_{ID} */
 $Psw \leftarrow \text{password}$

2. Message Sending: During message sending the following steps are followed by sender:

(a) $DID \leftarrow \text{Destination user name}$

/* sender filled the destination Username to which he wants to send the message that is stored as destination id DID*/

(b) $Msg \leftarrow \text{Plain text}$ /* sender write the message in to the plain text form*/

And send to the receiver

3. Storage of encrypted data in to the cloud server database : After sending the data the message archiver present into the cloud server data base performs the following actions:

(a) Packet formation :

(i) Message archiver receives the message and store it as one by one bit or character in message field of temporary created table

(ii) For each bit counter $\leftarrow \text{counter}+1$

(b) Encryption: message archiver perform the hash function on message field

$EK_{MSG} \leftarrow \text{Hash(message field)}$

(c) Storage of Master table: message archiver stores the master table into the cloud database which involves following fields:

Master Table $\leftarrow (S_{ID}, D_{ID}, EK_{MSG}, \text{Sequence no/counter})$

4. The master table is sent to the receiver via communication link.

5. End of algorithm

The brief description of algorithm is as follow



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Steps For sending encrypted data

Module 1: Authentication: The sender who wants to send the message firstly login at cloud server for authentication purpose where the sender username is taken as sender id.

Module 2: Message Sending: After login process the CSP requests to the sender to fill the information of destination identity and the message in the form of plain text that he wants to send to the receiver. After providing the required information to CSP the sender performs the action of sending the message by clicking on send button.

Module 3: Packet formation and encryption of message: Before sending the message from sender to receiver the CSP firstly converted the messages into packets and stores it as one by one bit. The counter is also incremented for each character/ bit of message which also deals with the problem of the collision of similar words caused due to the repetition of similar words during the transmission of long message. Cloud service provider also performs the hash function on message for encryption and maintaining their confidentiality.

Module 4: Storage of data as Master Table: The message archiver maintains the master table and stores it into Cloud database that has sender id, destination id, encrypted message and sequence number of whole message in its fields. The master table is then sent to the receiver through any of the network mediums.

The functional flow diagram and algorithm steps are described in figure 8 and 9 respectively.



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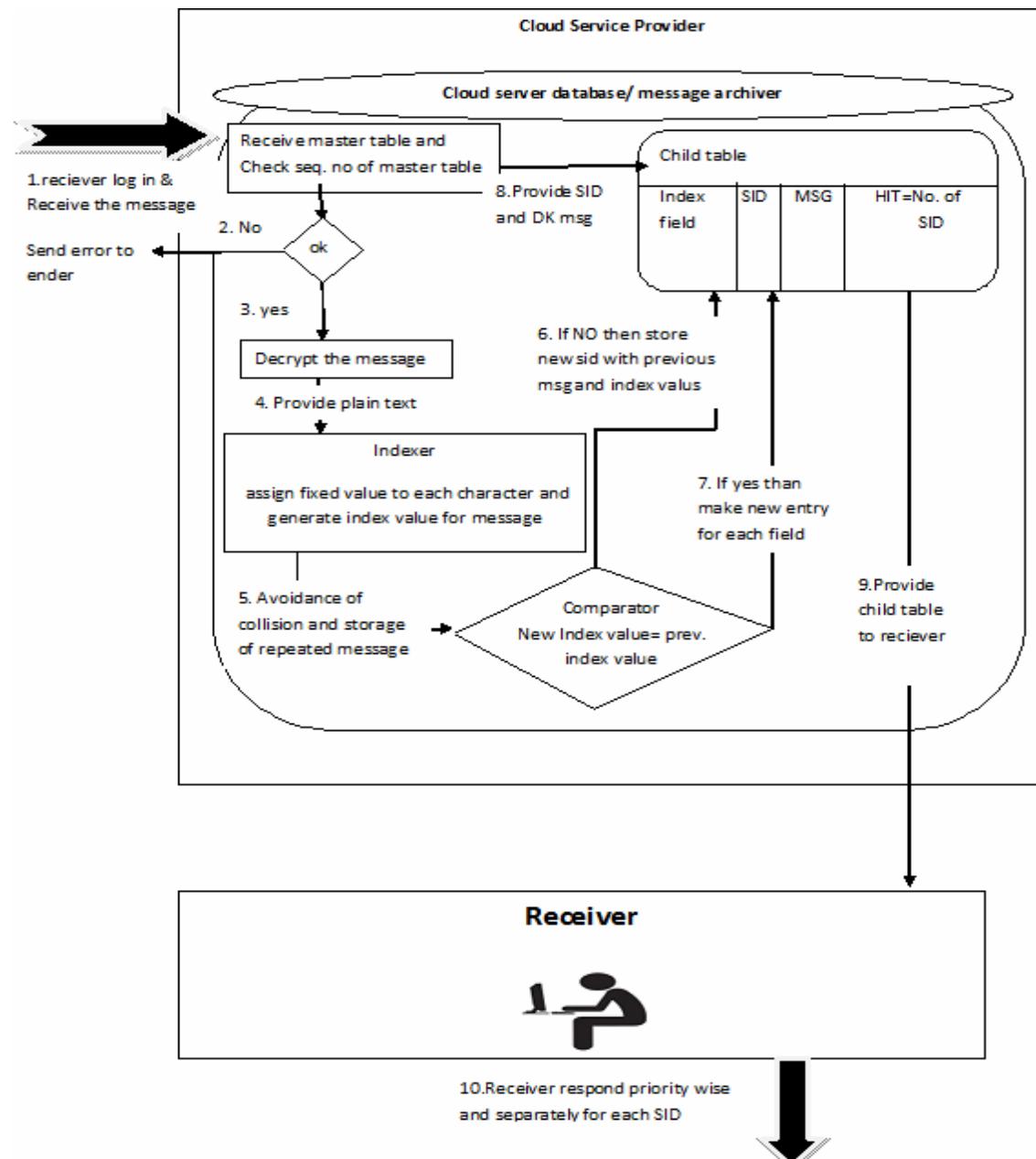


Figure 9. Functional Flow diagram for decrypting the message at receiver side



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ALGORITHM FOR DECRYPTION OF MESSAGE AND COLLISION AVOIDANCE AT RECEIVER SIDE

1. Receiving and decryption of incoming message
 - (a) Receiver side message archiver receives the master table
 - (b) Check for sequence number

If
sequence number is ok

Then
go to step 2

otherwise
send an error message to sender
2. Decrypt the message: $DK_{MSG} \leftarrow \text{Hash}_{\text{REVERSE}}(EK_{MSG})$
3. Creation of child table into the cloud database
 - (a) Assigned the fixed value for each character of decrypting message i.e.
 $FX_i \leftarrow i$ /* for each character I the fixed value is FX_i where $i = 1$ to n */
 - (b) Index value $\leftarrow (FX_1, FX_2, FX_3, \dots, FX_4)$
/* where the index value is equal to the concatenation of fixed values of all characters of decrypted msg */
 - (c) Message archiver stores the child table into the cloud database as follow:
Child table $\leftarrow (\text{index value}, S_{ID}, DK_{MSG}, HITS)$
4. Avoidance of collision and storage of repeated message:
 - (a) For each incoming message

If
 $\text{Index value}_{\text{new}} = \text{Index value}_{\text{prv}}$



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```
/*where message archiver match the index value for new
decrypting message with the index values of previous
messages existed in child table */
```

Then

Store new SID with the previous matched index value and

HIT← HIT+1

Otherwise

make a new entry for each field of child table

5. Receiver's response to the sender:
 - (a) Receiver responds priority wise and separately for each S_{ID} existed into the child table where greater the number of hits represents the higher the priority of message.
6. End of algorithm.

Steps for decrypting the message and avoidance of storage of repeated data at receiver side

Module 1: Check the sequence number for message: receiver side buffer or message archiver receives the Master table and check the sequence number of message. If the sequence number is correct then message archiver decrypts the encrypted message but if the sequence number is not correct then it sends the error message to sender.



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Module 2: creation of child table: The message archiver assigns the fixed value for each character of decrypted message i.e. assign FX_i for each i where $i=1$ to n . Message archiver also creates the index value for each message which is the concatenation of the fixed values that are assigned to each character of message and then store the child table into the cloud server that has following fields: index value for each message sender ids, decrypted message and hits that is incremented every time when a sender id is stored in child table.

Module 3: Avoidance of collision (loss of information) and storage of repeated messages: for each incoming message , message archiver always checks weather the index value of new arrived message is equal to the index value of any previously existed message. If it matches then only the sender id is stored in child table with the previously existed index value which prevents from the storage of same repeated message again in child table. But if the index values are not matched then message archiver saves a new entry for each field of child table. This technique also prevents from the loss of information due to the collision of same message at times because it always saves the sender id and also increment the number of hits whenever a new message arrived at receiver side.

Module 4: Priority wise response of receiver: The child table is sent to the receiver and receiver responds priority wise and separately for each sender id present into the child table. Where greater the number of hits represents highest the priority of message.

SIMULATION RESULTS AND ANALYSIS

SIMULATION AND RESULTS



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Secured Login

Username

Password

Simulation Performed and Tested by

Harpreet Aneja, M.Tech. Research Scholar, Lovely Professional University, Jalandhar, Punjab, India



Register / Sign Up

Name

Username

Password



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Welcome User :

Select Recipient	Name : Harpreet User ID : 1 Username : harpreet
Message Panel	
<input type="button" value="Send Message"/>	

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Secured Login

Username	<input type="text" value="harpreet"/>
Password	<input type="password" value="*****"/>
<input type="button" value="Submit"/>	

Simulation Performed and Tested by

Harpreeet Aneja, M.Tech. Research Scholar, Lovely Professional University, Jallandhar, Punjab, India



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Screenshot of a Firefox browser window showing an email inbox from "magmaconsultancy.com/harpreetthesis". The inbox contains one message with the following details:

Recipient ID 1
 Message Hello
 Sender ID 1
 Query Execution Time 0.00018191337585449
 Date and Time June 7, 2012, 3:08 pm

Below the inbox, a table titled "Fragmentation of Entire Message for Avoidance and Removal of Collision and Redundancy" is displayed:

Sr. N.	Character	Character Index	MD5 Hash	SHA1 Hash	Base 64 Encoding
1	H	0	c1d9f50f86825a1a2302ec2449c17196	7cf184f4c67ad58283ecb19349720b0cae756829	SA==
2	e	1	e1671797c52e15f763380b45e841ec32	58e6b3a414a1e090dfc6029add0f3555ccbba127f	ZQ==
3	I	2	2db95e8e1a9267b7a1188556b2013b33	07c342be6e560e7f43842e2e21b774e61d85f047	bA==
4	I	3	2db95e8e1a9267b7a1188556b2013b33	07c342be6e560e7f43842e2e21b774e61d85f047	bA==
5	o	4	d95679752134a2d9eb61dbd7b91c4bcc	7a81af3e591ac713f81ea1efe93dcf36157d8376	bw==
6		5	7215ee9c7d9dc229d2921a40e899ec5f	b858cb282617fb0956d960215c8e84d1ccf909c6	IA==
7	H	6	c1d9f50f86825a1a2302ec2449c17196	7cf184f4c67ad58283ecb19349720b0cae756829	SA==



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Sr. N.	Sender ID	Receiver ID	Message Transmitted	MD5 Hash Encryption	SHA1 Hash	Base 64 Encoding	Message Transmission Time (in Microseconds)	Date and Time
1	3	1	adasadg	795571285282b603cd789b66b5421179	59d001b1dca82f5c278bd6f81c683ce9c92c6181	c2Rhc2FkZn==	0.0012178421020508	May 25, 2012, 8:07 am
2	3	2		d41d8cd98f0b204e9800998ecf3427e	da39a3ee5e6b4b0d3255bfe95601890af80709		0.0025181770324707	May 25, 2012, 8:42 am
3	3	1		d41d8cd98f0b204e9800998ecf3427e	da39a3ee5e6b4b0d3255bfe95601890af80709		0.0012450218200684	May 25, 2012, 8:42 am
4	3	1	hello	5d41402abc4b2a76b9719d911017c592	aaf4c61ddcc5e8a2dabede0f3b482cd9aea9434d	aGVsbG8=	0.001331090927124	May 25, 2012, 8:42 am
5	3	1	hello	5d41402abc4b2a76b9719d911017c592	aaf4c61ddcc5e8a2dabede0f3b482cd9aea9434d	aGVsbG8=	0.001331090927124	May 25, 2012, 8:42 am
6	3	1	hello	5d41402abc4b2a76b9719d911017c592	aaf4c61ddcc5e8a2dabede0f3b482cd9aea9434d	aGVsbG8=	0.001331090927124	May 25, 2012, 8:42 am
7	3	1	hello	5d41402abc4b2a76b9719d911017c592	aaf4c61ddcc5e8a2dabede0f3b482cd9aea9434d	aGVsbG8=	0.001331090927124	May 25, 2012, 8:42 am
8	3	1	hello	5d41402abc4b2a76b9719d911017c592	aaf4c61ddcc5e8a2dabede0f3b482cd9aea9434d	aGVsbG8=	0.001331090927124	May 25, 2012, 8:42 am
9	3	1	hello	5d41402abc4b2a76b9719d911017c592	aaf4c61ddcc5e8a2dabede0f3b482cd9aea9434d	aGVsbG8=	0.001331090927124	May 25, 2012, 8:42 am
10	3	1	hello	5d41402abc4b2a76b9719d911017c592	aaf4c61ddcc5e8a2dabede0f3b482cd9aea9434d	aGVsbG8=	0.001331090927124	May 25, 2012, 8:42 am
11	3	1	hello	5d41402abc4b2a76b9719d911017c592	aaf4c61ddcc5e8a2dabede0f3b482cd9aea9434d	aGVsbG8=	0.001331090927124	May 25, 2012, 8:42 am



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User Inbox

Sr. N.	Sender ID	Sender Name	Message	Date and Time
1	3	Renu	sdasadg	May 25, 2012, 8:07 am
2	3	Renu		May 25, 2012, 8:42 am
3	3	Renu	hello	May 25, 2012, 8:42 am
4	3	Renu	hello	May 25, 2012, 8:42 am
5	3	Renu	hello	May 25, 2012, 8:42 am
6	3	Renu	hello	May 25, 2012, 8:42 am
7	3	Renu	hello	May 25, 2012, 8:42 am
8	3	Renu	hello	May 25, 2012, 8:42 am
9	3	Renu	hello	May 25, 2012, 8:42 am

Sr. N.	Sender ID	Sender Name	Message	Date and Time
1	3	Renu	sdasadg	May 25, 2012, 8:07 am
2	3	Renu		May 25, 2012, 8:42 am



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3	3	Renu	hello	May 25, 2012, 8:42 am
4	3	Renu	hello	May 25, 2012, 8:42 am
5	3	Renu	hello	May 25, 2012, 8:42 am
6	3	Renu	hello	May 25, 2012, 8:42 am
7	3	Renu	hello	May 25, 2012, 8:42 am
8	3	Renu	hello	May 25, 2012, 8:42 am
9	3	Renu	hello	May 25, 2012, 8:42 am



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10	3	Renu	hello	May 25, 2012, 8:42 am
11	3	Renu	hello	May 25, 2012, 8:42 am
12	3	Renu	dagsdagsdag	May 25, 2012, 9:00 am
13	3	Renu	dagsdagsdag	May 25, 2012, 9:00 am
14	3	Renu	dagsdagsdag	May 25, 2012, 9:00 am
15	3	Renu	dagsdagsdag	May 25, 2012, 9:00 am
16	3	Renu	dagsdagsdag	May 25, 2012, 9:00 am



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17	3	Renu	sagdsdag sdga gsda	May 25, 2012, 9:05 am
18	3	Renu	sagdsdag sdga gsda	May 25, 2012, 9:05 am
19	3	Renu	sagdsdag sdga gsda	May 25, 2012, 9:05 am
20	3	Renu	sagdsdag sdga gsda	May 25, 2012, 9:05 am
21	3	Renu	sagdsdag sdga gsda	May 25, 2012, 9:05 am
22	3	Renu	sagdsdag sdga gsda	May 25, 2012, 9:05 am
23	3	Renu	sagdsdag sdga gsda	May 25, 2012, 9:05 am



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24			sgsdag	May 26, 2012, 12:36 pm
25	3	Renu	hi how r u	May 27, 2012, 6:32 pm
26	2	Pooja	hello how r u	May 27, 2012, 6:39 pm
27	2	Pooja	hello how r u	May 27, 2012, 6:45 pm
28			hello h r u	May 27, 2012, 6:48 pm
29	3	Renu	hello how r u	May 28, 2012, 3:06 pm
30	1	Harpreet	Hello How r u ? Good Afternoon	May 30, 2012, 4:09 pm



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31	1	Harpreet	Hello	June 7, 2012, 3:08 pm
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SIMULATION

Sender ID	Recei ver ID	MD5 Hash Encryption	SHA1 Hash	Base 64 Encoding	Message Trans missio n Time (In Micros econds)
3	1	7955712852 82b603cdf8 9b66b54211 79	59d001b1dca8 2f5c278bd6f81 c683ce9c92c61 81	c2Rhc2FkZw==	0.0012 17842 10205 08
3	2	d41d8cd98f 00b204e980 0998ecf842 7e	da39a3ee5e6b 4b0d3255bfef9 5601890af807 09		0.0025 18177 03247 07
3	1	d41d8cd98f 00b204e980 0998ecf842 7e	da39a3ee5e6b 4b0d3255bfef9 5601890af807 09		0.0012 45021 82006 84
3	1	5d41402abc 4b2a76b971 9d911017c5 92	aaf4c61ddcc5e 8a2dabede0f3b 482cd9aea943 4d	aGVsbG8=	0.0013 31090 92712 4
3	1	5d41402abc 4b2a76b971 9d911017c5 92	aaf4c61ddcc5e 8a2dabede0f3b 482cd9aea943 4d	aGVsbG8=	0.0013 31090 92712 4
3	1	5d41402abc 4b2a76b971	aaf4c61ddcc5e 8a2dabede0f3b	aGVsbG8=	0.0013 31090



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		9d911017c5 92	482cd9aea943 4d		92712 4
3	1	5d41402abc 4b2a76b971 9d911017c5 92	aaf4c61ddcc5e 8a2dabede0f3b 482cd9aea943 4d	aGVsbG8=	0.0013 31090 92712 4
3	1	5d41402abc 4b2a76b971 9d911017c5 92	aaf4c61ddcc5e 8a2dabede0f3b 482cd9aea943 4d	aGVsbG8=	0.0013 31090 92712 4
3	1	5d41402abc 4b2a76b971 9d911017c5 92	aaf4c61ddcc5e 8a2dabede0f3b 482cd9aea943 4d	aGVsbG8=	0.0013 31090 92712 4
3	1	5d41402abc 4b2a76b971 9d911017c5 92	aaf4c61ddcc5e 8a2dabede0f3b 482cd9aea943 4d	aGVsbG8=	0.0013 31090 92712 4
3	1	5d41402abc 4b2a76b971 9d911017c5 92	aaf4c61ddcc5e 8a2dabede0f3b 482cd9aea943 4d	aGVsbG8=	0.0013 31090 92712 4
3	1	5d41402abc 4b2a76b971 9d911017c5 92	aaf4c61ddcc5e 8a2dabede0f3b 482cd9aea943 4d	aGVsbG8=	0.0013 31090 92712 4
3	1	ed1034c02f 9256384931 68d476fb1 60	bfdb730bf032fa 0d7892407f690 4c6e35dc8e74 0	ZGFnc2RhZ3NkYWc=	0.0011 86132 43103 03
3	1	ed1034c02f 9256384931 68d476fb1 60	bfdb730bf032fa 0d7892407f690 4c6e35dc8e74 0	ZGFnc2RhZ3NkYWc=	0.0011 86132 43103 03
3	1	ed1034c02f 9256384931 68d476fb1 60	bfdb730bf032fa 0d7892407f690 4c6e35dc8e74 0	ZGFnc2RhZ3NkYWc=	0.0011 86132 43103 03



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3	1	ed1034c02f 9256384931 68d476fb1 60	bfdb730bf032fa 0d7892407f690 4c6e35dc8e74 0	ZGFnc2RhZ3NkYWc=	0.0011 86132 43103 03
3	1	ed1034c02f 9256384931 68d476fb1 60	bfdb730bf032fa 0d7892407f690 4c6e35dc8e74 0	ZGFnc2RhZ3NkYWc=	0.0011 86132 43103 03
3	1	e2a17c6b79 4245417a8e 4dcbeae68a e3	07475db5a0d4 139dd9e916d5 ac9f5fc503ddc2 82	c2FnZHNkYWcNCnNkZ2ENCmdzZGE=	0.0016 72029 49523 93
3	1	e2a17c6b79 4245417a8e 4dcbeae68a e3	07475db5a0d4 139dd9e916d5 ac9f5fc503ddc2 82	c2FnZHNkYWcNCnNkZ2ENCmdzZGE=	0.0016 72029 49523 93
3	1	e2a17c6b79 4245417a8e 4dcbeae68a e3	07475db5a0d4 139dd9e916d5 ac9f5fc503ddc2 82	c2FnZHNkYWcNCnNkZ2ENCmdzZGE=	0.0016 72029 49523 93
3	1	e2a17c6b79 4245417a8e 4dcbeae68a e3	07475db5a0d4 139dd9e916d5 ac9f5fc503ddc2 82	c2FnZHNkYWcNCnNkZ2ENCmdzZGE=	0.0016 72029 49523 93
3	1	e2a17c6b79 4245417a8e 4dcbeae68a e3	07475db5a0d4 139dd9e916d5 ac9f5fc503ddc2 82	c2FnZHNkYWcNCnNkZ2ENCmdzZGE=	0.0016 72029 49523 93
3	1	e2a17c6b79 4245417a8e 4dcbeae68a e3	07475db5a0d4 139dd9e916d5 ac9f5fc503ddc2 82	c2FnZHNkYWcNCnNkZ2ENCmdzZGE=	0.0016 72029 49523 93
3	1	e2a17c6b79 4245417a8e 4dcbeae68a e3	07475db5a0d4 139dd9e916d5 ac9f5fc503ddc2 82	c2FnZHNkYWcNCnNkZ2ENCmdzZGE=	0.0016 72029 49523 93
3	2	bb3a85aca6 32e03955f8 9c3f0dc3fa9	656d0201d0e1 db671c76735e 04e3ec49b2f93	YXNkc2RhZw==	0.0019 37866 21093



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3	2	bb3a85aca6 32e03955f8 9c3f0dc3fa9 8	656d0201d0e1 db671c76735e 04e3ec49b2f93 378	YXNkc2RhZw==	0.0019 37866 21093 75
3	2	bb3a85aca6 32e03955f8 9c3f0dc3fa9 8	656d0201d0e1 db671c76735e 04e3ec49b2f93 378	YXNkc2RhZw==	0.0019 37866 21093 75
3	2	bb3a85aca6 32e03955f8 9c3f0dc3fa9 8	656d0201d0e1 db671c76735e 04e3ec49b2f93 378	YXNkc2RhZw==	0.0019 37866 21093 75
3	2	bb3a85aca6 32e03955f8 9c3f0dc3fa9 8	656d0201d0e1 db671c76735e 04e3ec49b2f93 378	YXNkc2RhZw==	0.0019 37866 21093 75
3	2	bb3a85aca6 32e03955f8 9c3f0dc3fa9 8	656d0201d0e1 db671c76735e 04e3ec49b2f93 378	YXNkc2RhZw==	0.0019 37866 21093 75
1	3	8b1a9953c4 611296a827 abf8c47804 d7	f7ff9e8b7bb2e0 9b70935a5d78 5e0cc5d9d0abf 0	SGVsbG8=	0.0010 02788 54370 12
1	4	c92a8c4c96 39c8eea645 ae900d2d55 d0	974d929ed4b5 aab1d1fcf02fe a21a48eafc6c4 8	aGVsbG8gdw==	0.6835 40105 8197
1	6	26efe43225 6ef88b79f75 59130a982b 5	3ebe997f741f1 4c09be65dff88 2610f851e3f66 2	Z2pmZ2oNCmprbGsNCm4sbW5rLA0KamJtYiw NCmtubW4gDQoNCm5odm1iDQpibm5iDQpub WJtYg0Kbm1iLG1uaywNCg==	0.0009 73939 89562 988
	1	59cab9ac39 de16f14029 12e651b4a6 f2	ee13de73a30e ecb7cf7d4c161 d6a4f4299b1f8 57	c2dzZGFn	0.9271 28076 55334



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3	4	163745a94a d2451e897c d2d78723e8 1b	a792ab289441 80df2a4d9bb46 86c6ae08592d a35	aGVsbG8NCmhvdyBydSA/	0.0519 76919 17419 4
3	3	26407e224d 07f274d9f49 53d79a616e c	eae8644396f12 24d065b69050 6de3525c25a7f 68	ZGdzYQ0KaGVsbG8=	0.0003 76939 77355 957
3	1	b68682e147 64463aabec 40f464a9c9 98	00db1a0b08b2 8e8541a6b4bb d8afab09544c7 ad0	aGkgaG93IHlgdQ==	0.0005 92947 00622 559
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2	1	a71698d590 eb6791392f 770a45b510 73	d4d1a60f8d365 a55c8a3abf406 5c09ec6409def 0	aGVsbG8gaG93IHlgdQ==	0.0002 45809 55505 371
	1	08f50adbe7 75ce25c981 76453d2c78 e3	1b62a81b7877f 0613a92fd68c2 9fc75e6a2fb6c 2	aGVsbG8gaCByIHU=	0.0002 42948 53210 449
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3	3	f582339ce7 5b24996e5e add71b3168 51	2fa9aaa270c16 9743dd17c3bb a6cbfd0f605f06 a	aGVsbG9ob3cgciB1Pw0K	0.0005 22136 68823 242
1	1	15b061bfdf9 8b81cee651 c2b1d98e3b 1	47ca001c1bf9d e35181bfc6069 85aa6c131d17 e3	SGVsbG8gSG93IHlgdSA/DQoNCkdvb2QgQW Z0ZXJub29u	0.0005 70058 82263 184
1	1	8b1a9953c4 611296a827 abf8c47804	f7ff9e8b7bb2e0 9b70935a5d78 5e0cc5d9d0abf	SGVsbG8=	0.0001 81913 37585



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		d7	0		449
8	1	683aef499f8 2c9ea93624 e9f4117a25 4	777b73064a6b 6a1c981276ad 2a35c5b68b75f 817	aGVsbG8NCmdvb2QgbW9ybmluZw==	0.0637 18080 52063
8	2	192ceeb80e 05b8f14514 801ec1fb3f d	dd81b6bb040d 72b522af0fb0a ae18b2637ac4 c7b	aHNkYTtzZGFdDQpzZGENCg0Kc2Rhc2RhDQ pzZmRhDQpzZmdhDQoNCnNkDQoNCg0KDQ oNCg0KDQoNCg0Kc2FnYQ0KYXNkDQphc2R nDQo=	0.0004 63962 55493 164
9	2	f057f5ccb87 a310534b9d da3a69a72c 4	650f97e596e63 e92d9447619d 38e8a83e0bec 8b9	aGlp	0.0014 51015 47241 21

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